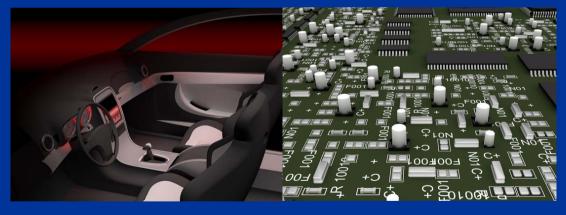
# Aperiodic traffic in response time analyses with adjustable safety level

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Complexity Mastered

#### Validation is a key activity in automotive systems design Personal view on the developments

Mostly ahead of us!

« correctness by construct » and optimal synthesis

Probabilistic analysis system level

« Worst-case » deterministic analysis system level

**Probabilistic analysis (sub-system)** 

« Worst-case » deterministic analysis (sub-system)

« Smart » monitoring tools

Simulation tools (co-simulation, HIL)

1994

1997

2009







#### Probabilistic analysis is needed

- Systems are not designed for the worst-case (provided it is rare enough!)
- Reliability/Safety are naturally expressed and assessed in terms of probability (e.g. < 10^-9 per hour)
- Deterministic assumptions are sometimes unrealistic or too pessimistic, e.g.:
  - Worst-Case Execution Time on modern platforms,
  - Aperiodic activities: ISR, frame reception,
  - **...**
- Faults/errors are not deterministic (and better modeled probabilistically)







### Accounting for the aperiodic traffic

- Transmission patterns can hardly be characterized: purely aperiodic, mixed periodic/aperiodic, etc
- Aperiodic frames do jeopardize RT constraints
- Few approaches in the litterature:
  - deterministic approaches, such as sporadic, generally lead to unusable results (e.g., ρ>1)
  - Average case probabilistic approach not suited to dependability-constrained systems
  - Probabilistic approaches with safety adjustable level, see paper ref[6] and ref[7]







#### Approach advocated here

- 1) Measurements / data cleaning
- 2) Modeling aperiodic traffic arrival process
- 3) Deriving aperiodic Work Arrival Process (WAF)
- 4) Integrating aperiodic WAF into schedulability analysis





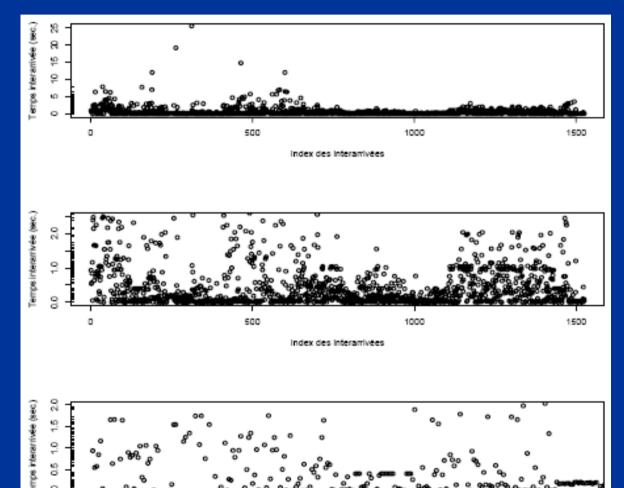


#### Data trace analysis

y: aperiodic interarrival times - x: index of interarrivals



Approximate because what is seen on the bus is not the actual arrival process at ECU level! can be handled



x: [0-1500]

y: [0-25]

x: [0-1500]

y: [0-2.5]

x: [0-250]

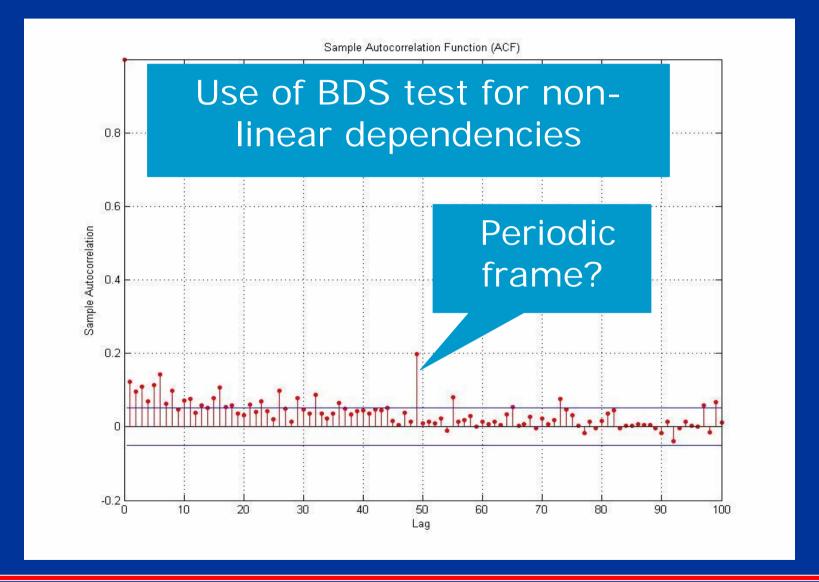
y: [0-2]







#### Question: are interarrival times i.i.d.?



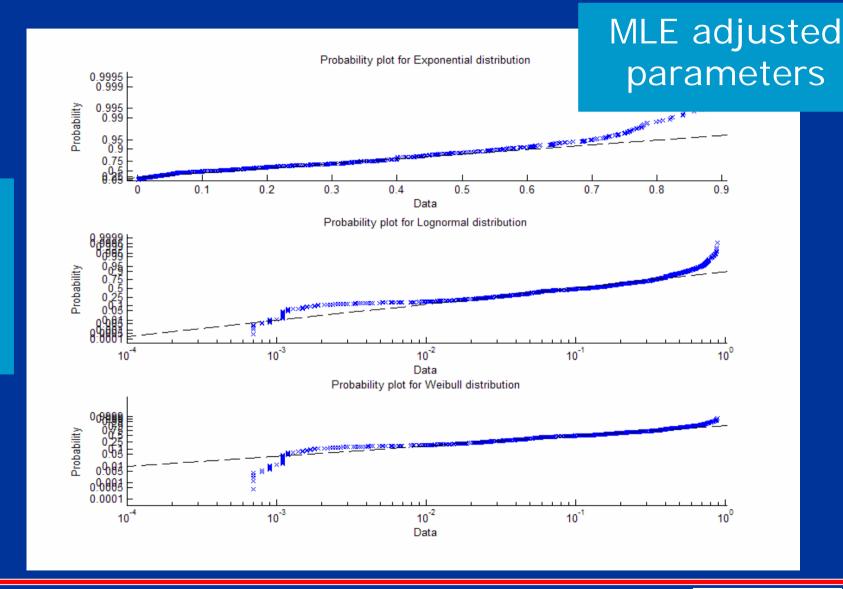






### Distribution fitting for aperiodic interarrival: 3 candidates here

Kolmo. Smi. and χ2 tests to confirm visual impression

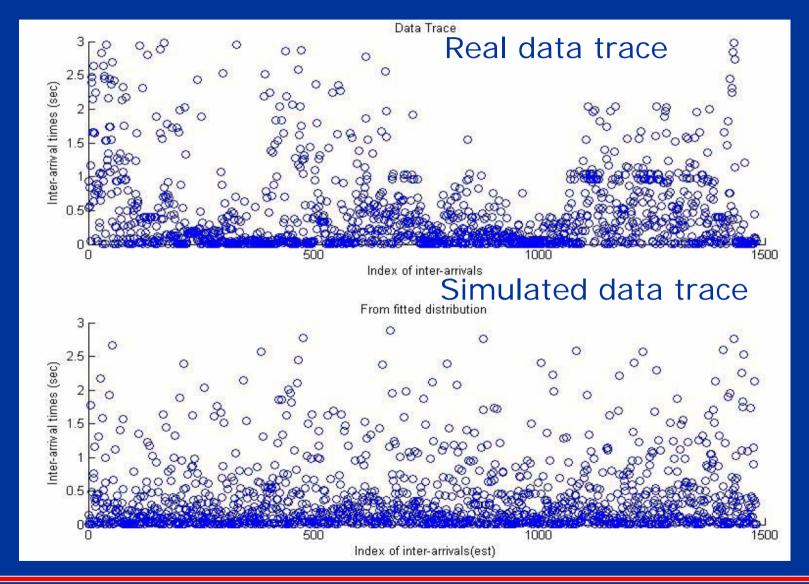








# Captured data trace VS random trace generated with MLE-fitted Weibull









### Deriving the aperiodic WAF

- S(t): aperiodic WAF
- X(t): stochastic process which counts the number of aperiodic frames in time interval t
- "smallest" S(t) such that the probability of X(t) being larger than or equal to S(t) is lower than a threshold  $\alpha$

$$\hat{S}(t) = \min\{S(t) \mid \Pr[X(t) \ge S(t)] \le \alpha\}$$

By simulation, numerical approximation or analysis (simplest cases such as exp.)

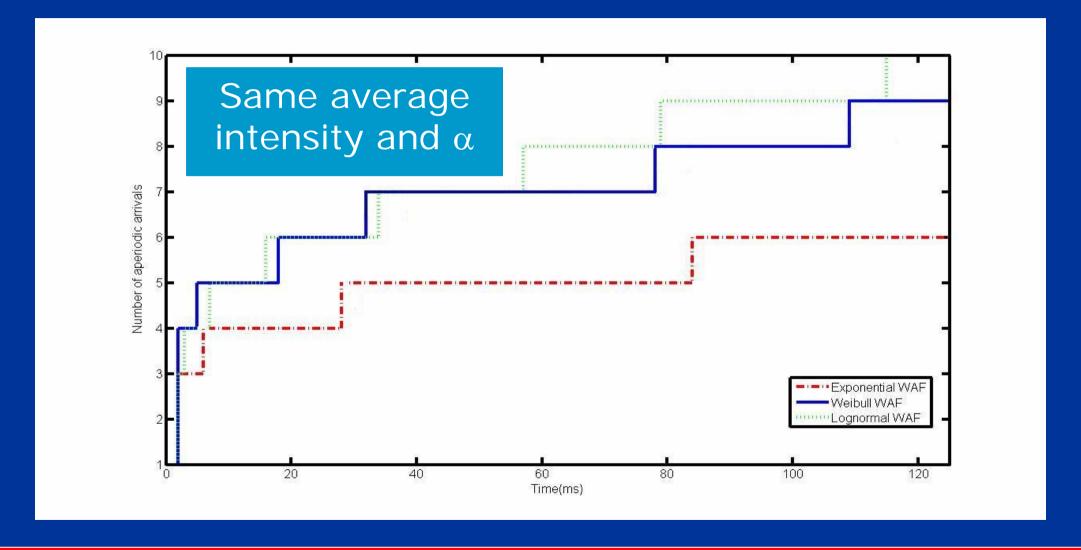
Design choice: e.g., 10<sup>-9</sup>







### Aperiodic WAF depends on the underlying interarrival distribution









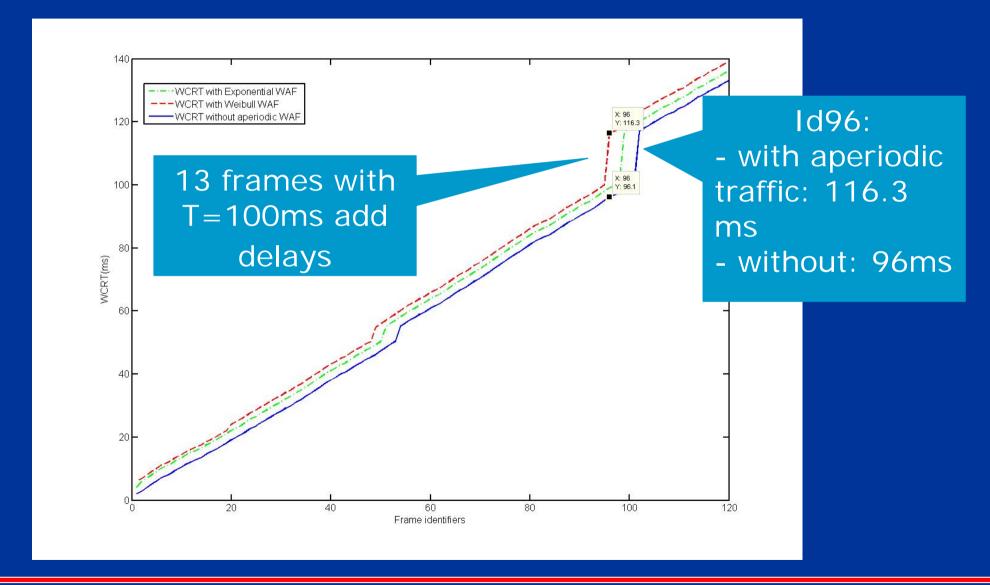
#### Case-study on a typical body network

- Body network benchmark generated using GPLlicensed Netcarbench
- Characteristics:
  - 125kbps, 16 ECUs, 105 CAN frames with deadlines equal to periods and 1 to 8 bytes of data.
  - Total periodic load is equal to 41%
- NETCAR-Analyzer for WCRT computation
- 3% aperiodic traffic
- 7 byte aperiodic frames
- $\alpha = 10^{-4}$





# Worst-case response times with/out aperiodic traffic (3%)

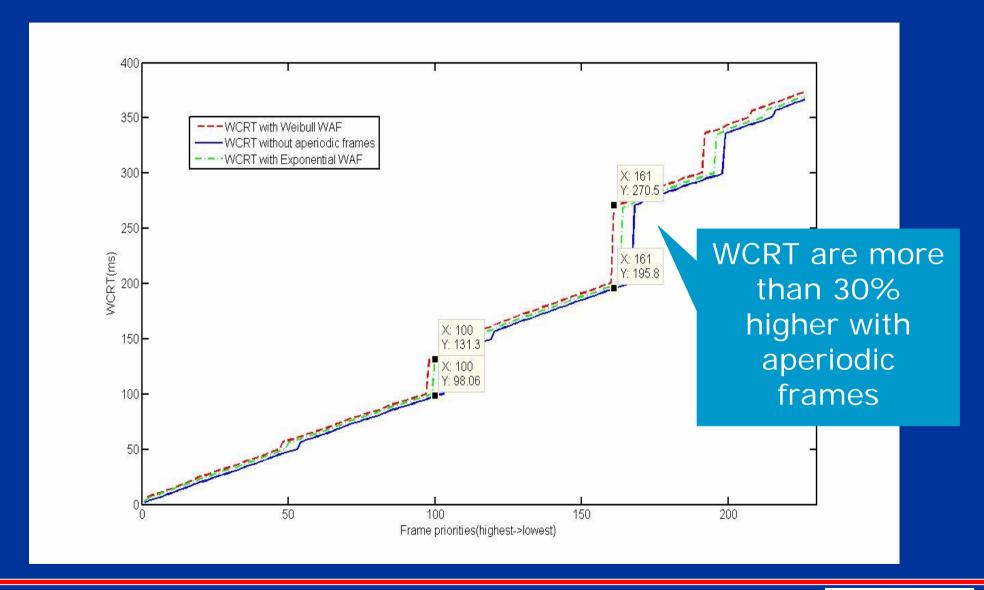








#### On a more loaded network...









#### Observations

- In this context where the periodic load is relatively small and the aperiodic traffic is limited (3%) one observes:
- aperiodic traffic significantly impacts the worst-case response times of the periodic frames (more than +30% sometimes).
- the exact model of the aperiodic traffic plays some role
- depends on the priority of the aperiodic frames (working on this)
- Measured arrival time on bus at which the frames started to be transmitted can be different than time at which the transmission requests were issued





#### Conclusion

- Chosen dependability requirements are met while pessimism kept to minimum:
  - Practical approach
  - Real data are required
  - Can be extended to the non i.i.d. case (not needed here)
- What is needed now is a system level approach that
  - Can handle arbitrary activation processes
  - goes beyond the i.i.d. case (for dependability assessment)







#### References







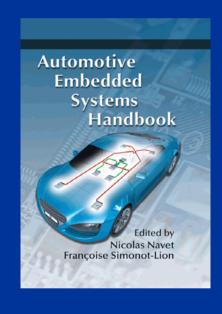
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- [2] P. Wallin, Axelsson, A Case Study of Issues Related to Automotive E/E System Architecture Development, IEEE International Conference and Workshop on the Engineering of Computer Based Systems, 2008.
- [3] H. Hansson, M. Nolin, T. Nolte, Beating the Automotive Code Complexity Challenge: Components, Models and Tools, National Workshop on High-Confidence Automotive Cyber-Physical Systems, 2008.

#### Dependability / probabilistic framework

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- [5] D. Khan, N. Navet, B. Bavoux, J. Migge, "Aperiodic Traffic in Response Time Analyses with Adjustable Safety Level", IEEE ETFA2009, Mallorca, Spain, September 2009.
- [6] N. Navet, Y-Q. Song, F. Simonot, "Worst-Case Deadline Failure Probability in Real-Time Applications Distributed over CAN (Controller Area Network)", Journal of Systems Architecture, Elsevier Science, vol. 46, n°7, 2000.
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#### Questions / feedback?



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